

ARCURVE presents...

# Alberta Collegiate Programming Contest 2024

Hosted by:  **Kattis**

Run by:  **COMPETITIVE  
PROGRAMMING  
CLUB**



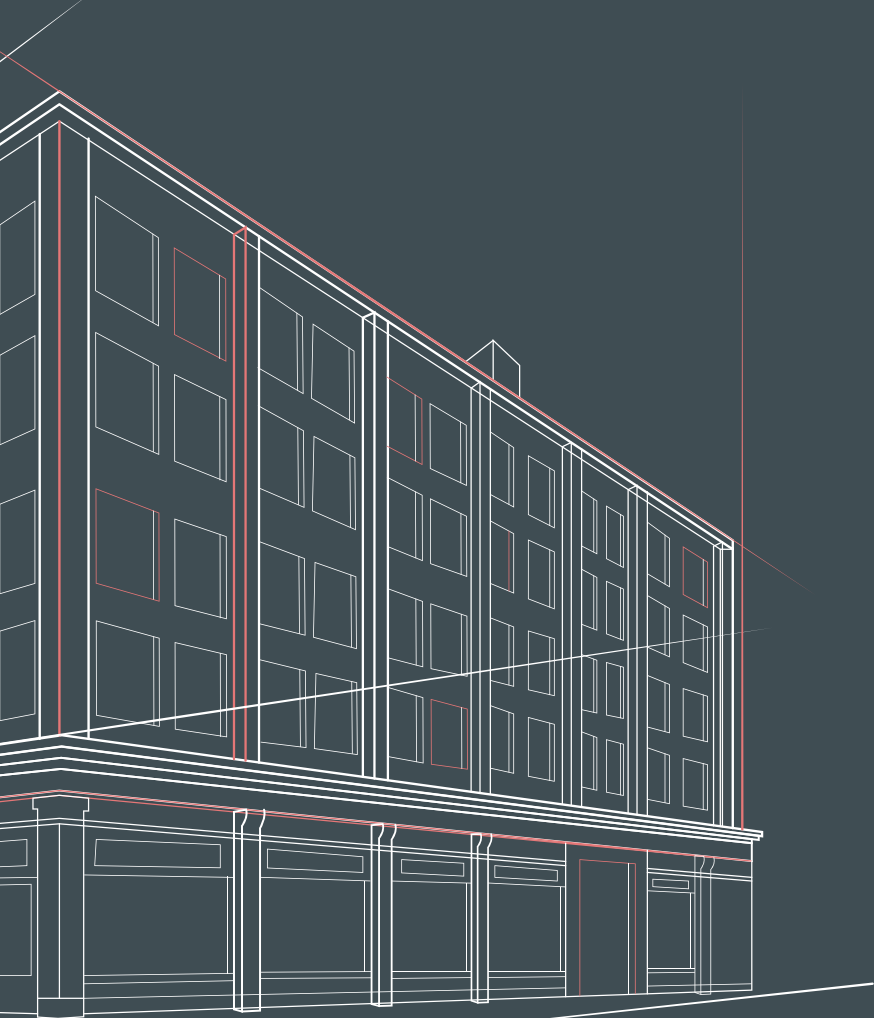
..... **Thank you!** .....

Big shoutout to our Problem  
setters, Kattis, and executive  
team for making this  
happen!

Problem setters: Taylor Wong, John Zheng, Tung Nguyen,  
Noah Weninger, Dante Bencivenga, Zachary Friggstad,  
Charlie Zheng, Ian DeHaan, and Howard Cheng

Kattis Representative: Fredrik Niemelä

Arcurve Representatives: Mike Bauer, Joel Pollard, and Rose  
Muhammed

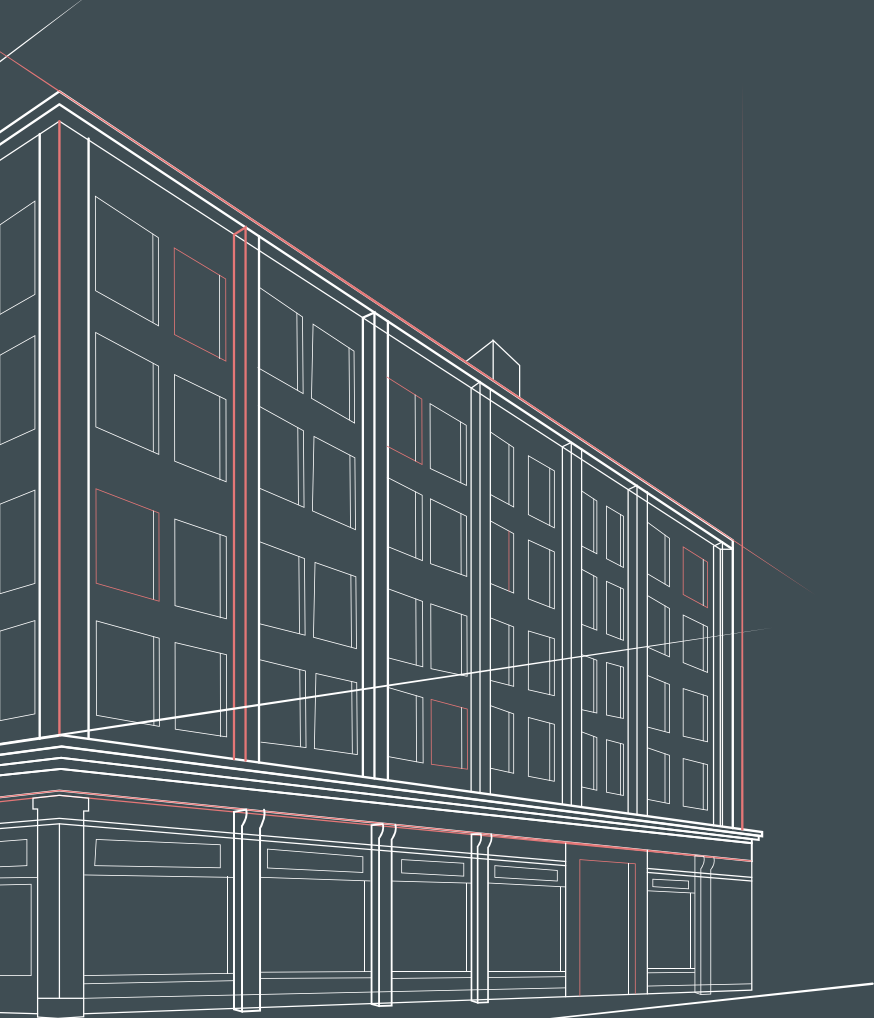


01

# Statistics



Lets see how the contest  
went!



# 02

## Problem Solutions

Recap of problems and  
brief solutions

# Snowfall

## by John Zheng

First solve Div 1: definitely not alberta sapphire or anything like that

### First solve Div 2: DTS

Predicted: 98% solved

Actual Div 1: 100%

Actual Div 2: 100%

- Read all the inputs, and based on the value of  $t$  (standing for type), increase or decrease the amount of accumulated snow (starting with no snow)
- Snow depth will not become negative, so keep it above 0

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## A Little to the Right by John Zheng

**First solve Div 1: but have you ever had a donut?**

**First solve Div 2: Git Good**

Predicted: 92% solved    Actual Div 1: 60.87%    Actual Div 2: 31.82%

- There are only  $p$  properties, so at most  $p$  different orderings
- Sort list  $p$  times, one time for each property
  - Make sure that property is strictly increasing (no duplicates)
  - Make sure this ordering isn't identical to a previous ordering
  - Input size is small enough for  $O(p^2 \cdot n)$

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# Melting Snow


## By Ian DeHaan

### First solve: massachusetts 4.0



Predicted: 100% solved

Actual Div 2: 70.72%



```
s, p = map(int, input().split())
p = float(p/100)

# ends when the gain of snow is equal to the
# amount disappearing
# i.e. when s = out*p
# out = s/p

print("%.6f" % float(s/p))
```

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# Handheld Fan

by Zachary Friggstad

## First solve: potato



Predicted:      Actual Div 2: 45.45%

Try all starting hours  $i$ . Step along hours  $i, i+1, i+2$ , etc. until the total requirement of the fan would exceed its capacity.

Output the longest answer you found over all possible starting times.

Running time:  $O(n^2)$

Can do in  $O(n)$  by not resetting the end interval of the scan when try the next  $i$  (just resume from there).



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# Alpine Agility by Taylor Wong

## First solve Div 1: Free Pizza



## First solve Div 2: Stack Underflow

● Predicted: 40% solved

Actual Div 1: 43.48% solved

Actual Div 2: 9.90% solved

Standard DP problem.

The fastest you can get to slope[i, j] (speed[i, j]) is

max(

speed[i - 1, j - 1] + h[i - 1, j - 1] - 150,

speed[i - 1, j] + h[i - 1, j] - 100,

speed[i, j - 1] + h[i, j - 1] - 100

) - h[i, j]

# Hockey Fans

## by Zachary Friggstad

First solve: definitely not alberta sapphire or anything like that

Predicted: Actual Div 1: 39.13% Actual Div 2: 9.09%

Binary search the answer. For each query  $D$ , find all consecutive intervals where the decibel level is  $\geq D$ . You can do  $\text{floor}(L/D)$  shouts in such an interval of length  $L$ . Summing over all intervals, can you do the required number of shouts?

Running time:  $O(n \log \text{MAXDECIBEL})$

Alternatively, sort the times in reverse order by their decibel level. In this sorted order, "activate" the interval and merge it with adjacent activated intervals. Update the running total of  $\text{floor}(L/D)$  over all active intervals. Stop when this is the number of shouts.

Running time:  $O(n \log n)$

# Topographic Isolation

## by John Zheng

First solve Div 1: definitely not alberta sapphire or anything like that

### First solve Div 2: Stack Underflow

Predicted: 30% solved      Actual Div 1: 21.74%      Actual Div 2: 18.18%

- Naive solution: For each peak, scan left until a equal or higher point, and then scan right for a equal or higher point
  - $O(n^2)$ , too slow
- Optimization: instead of scanning all the way left then all the way right, alternate between left and right and break immediately once equal or higher point is found
  - $O(n \log n)$ , proof is left as an exercise for the reader
- Can also use monotonic stack for  $O(n)$

# Four Questions

by Noah Wenginger and Ian Parrish

## First solve:

Predicted: 10% solved

Actual Div 1: 8.33%

- Start by querying 100,000 and suppose the answer is  $r$
- Then  $p$  must be a prime factor of  $100,000^2 - r$
- It can be shown that there are at most 8 distinct prime factors
- Binary search over products of the 8 prime factors to narrow it down. This requires  $\log(8)=3$  more queries.

# Snow Way Out

## by Dante Bencivenga

First solve Div 1: definitely not alberta sapphire or anything like that

First solve Div 2: Linux Ladies

Predicted: 50% solved

Actual Div 1: 17.39%

Actual Div 2: 9.09%

- Take a closer look at the formula for *squared* distance:  
 $(x-x')^2 + (y-y')^2$ 
  - If we only vary the x coordinate and keep y the same, only the x contribution to squared distance changes
  - Therefore, we can independently narrow down the x coordinate first while keeping y constant, and then narrow down y, keeping x constant
- We have reduced the problem to finding one coordinate from 0 to 9 using three "?" queries
- For one coordinate there are only 10 possibilities, so each search path can be hard-coded and used the same way for both x and y

# Uniform Variation by John Zheng

**First solve: definitely not alberta sapphire or anything like that**

Predicted: 50% solved

Actual Div 1: 13.04%

Actual Div 2: 0%

- Two cases:
  - If there's a single type of gear which has variation, there can only be up to four people
    - only four colours
  - To get five or more people in a photo, they must have identical gear
- All combinations of two people are valid
- Brute force all combinations of 3 and 4 people
- Check exact duplicates for 5 more more people

# Wrapping Snowballs by Charlie Zheng

First solve Div 1: LOSA Meow Meows ;3 in 47 minutes

First solve Div 2: Stack Underflow in 120 minutes

Predicted: 25% solved

Actual Div 1: 13.04%

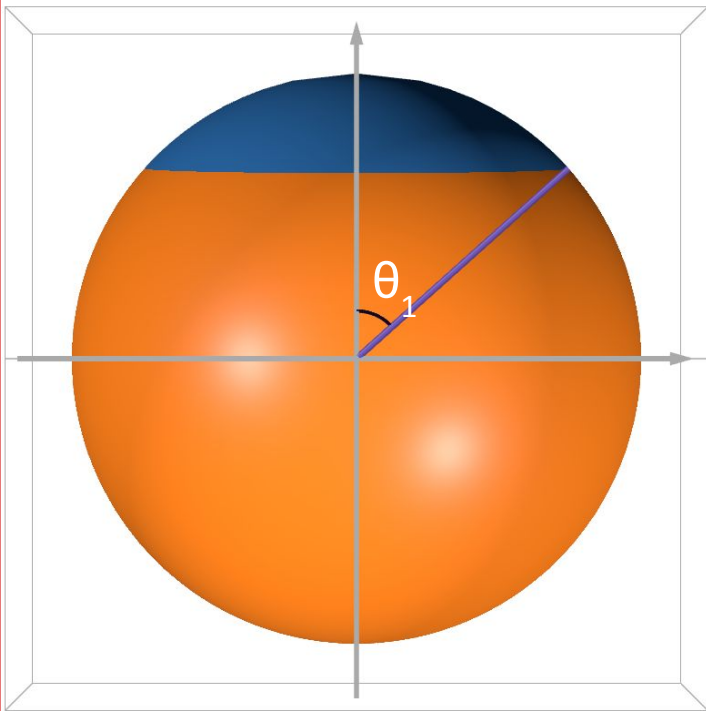
Actual Div 2: 9.09%

- Arc length of part of sphere that can be covered by the paper is  $R_{paper}$
- $\theta_1 = R_{paper} / R_{ball}$
- If the entire ball can be covered ( $\theta \geq 2\pi$ ), the answer is the SA of the sphere:  $4 \cdot \pi \cdot R_{ball} \cdot R_{ball}$
- 
- Otherwise, cylindrical integration to get surface area of the top of the sphere that can be covered:

$$SA = R_{ball} \int_0^{\theta_1} 2\pi R_{ball} \times \sin(\theta) d\theta$$

- You can compute this integral numerically or solve it and get:

$$SA = 2\pi \times R_{ball} \times R_{ball} \times (1 - \cos(\theta_1))$$



# Slippery Floor by Tung Nguyen

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First solve div 1: definitely not alberta sapphire or anything like that

First solve div 2: Linux Ladies



● Predicted: 70% solved

Actual Div 1: 13.04%

Actual Div 2: 4.55 %

- BFS on different stopping locations and four possible directions
  - Brute force scanning for the next obstacle is too slow
- Keep sorted list for every horizontal and vertical coordinate in the input
- Binary search on the sorted list to find next obstacles for pushing directions



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# Speedy Slopes by Taylor Wong

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**First solve: definitely not alberta sapphire or anything like that**

Predicted: 20% solved

Actual Div 1: 4.35% solved

SSSP from Northwest -> Southeast (Dijkstra's):

Vertices = state including the slope and **speed**.

Edges Weights = distance to next / new speed (time taken to move)

Optimization:

If you arrive at a slope with less speed then before you can skip it. Instead of tracking seen slopes, track the max speed you've reached each slope with.

# Thank you for participating!

If you haven't already join our Discord it's our main source of communication:

<https://discord.gg/7ZAqJvX>

Find the rest of our socials on our website:

<https://cpc.cpssc.ucalgary.ca/>

